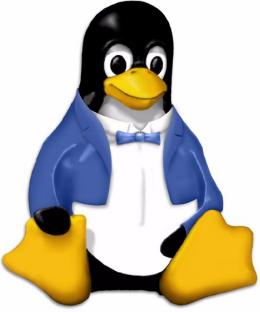




# Device-driven I/O for Implicit Paging Operations

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# Outline

- K42 overview
- The K42 I/O system
- The device-driven model
- Implementation & experimentation
- Further work



# K42 Overview

- Research OS
  - Intended for easy experimentation
- Microkernel-style architecture
- Fast user/kernel-space message passing system
  - Protected Procedure Call (PPC)
  - Stub compiler to easily define new calls
- Currently working on 64-bit POWER machines
- Compatible with Linux API & ABI



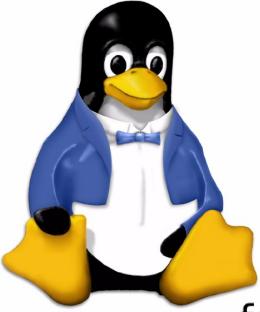
# How it works now...

- K42 I/O subsystem
  - File objects are kept in userspace
  - write() is performed asynchronously
    - essentially a memcpy()
- Walk through a fsync...

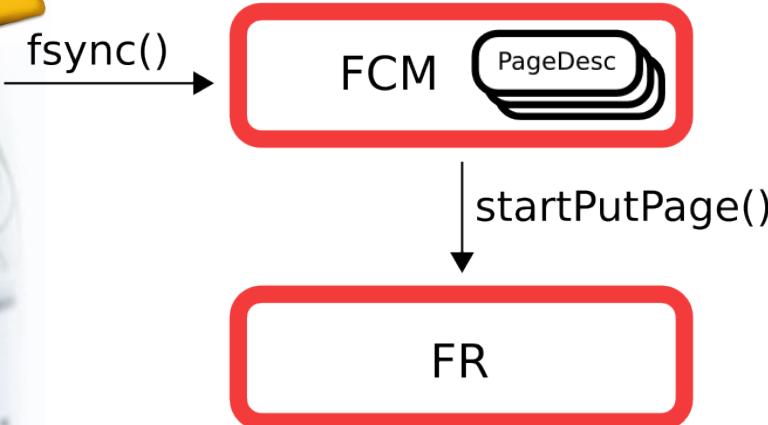
# Current fsync Path



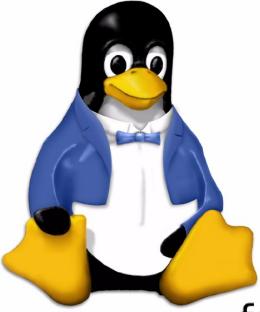
File Cache Manager receives fsync()



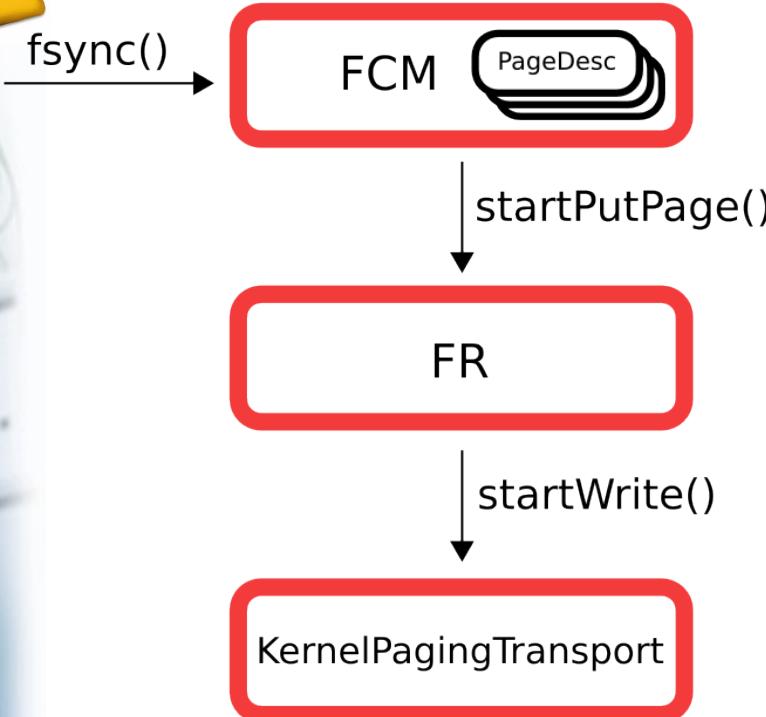
# Current fsync Path



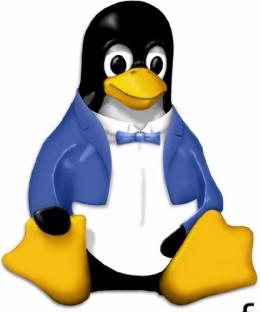
File Representative manages reference to userspace file



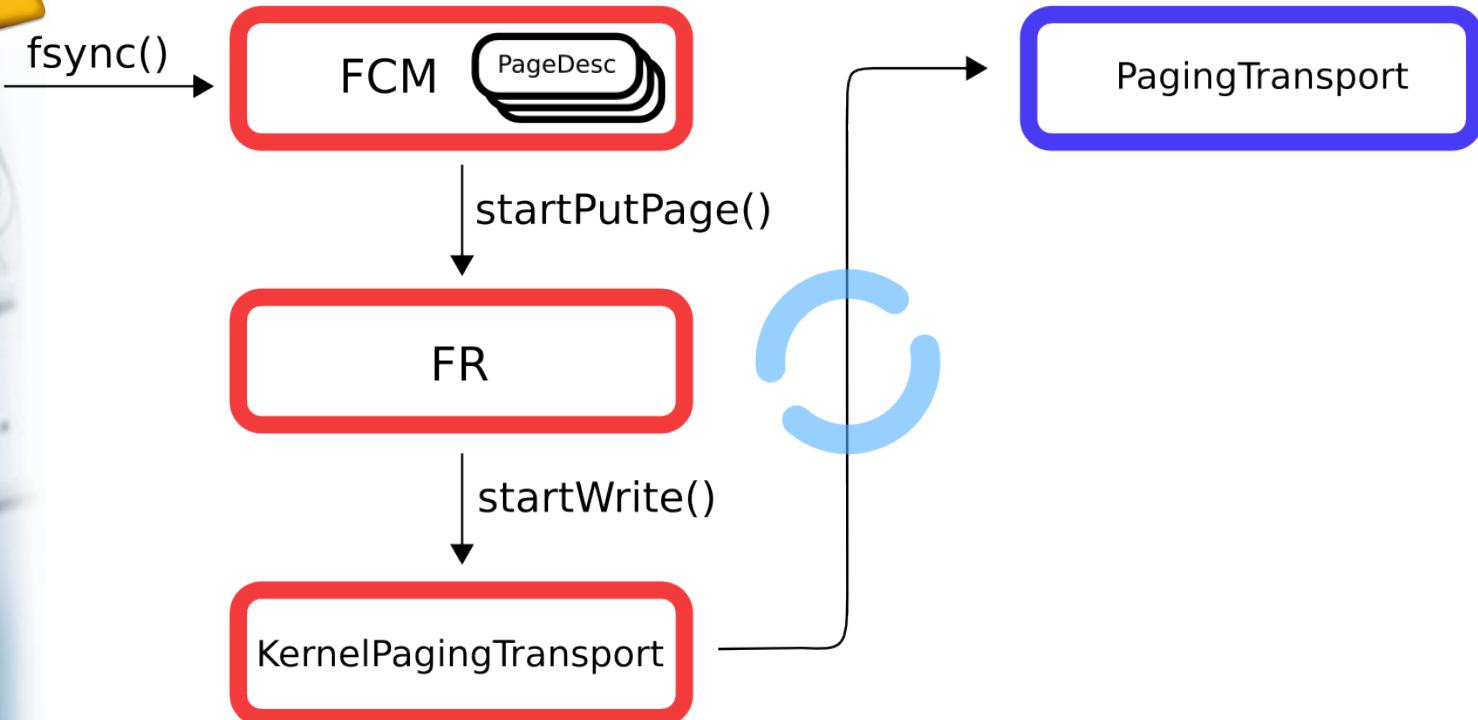
# Current fsync Path



KernelPagingTransport passes I/O request to userspace...



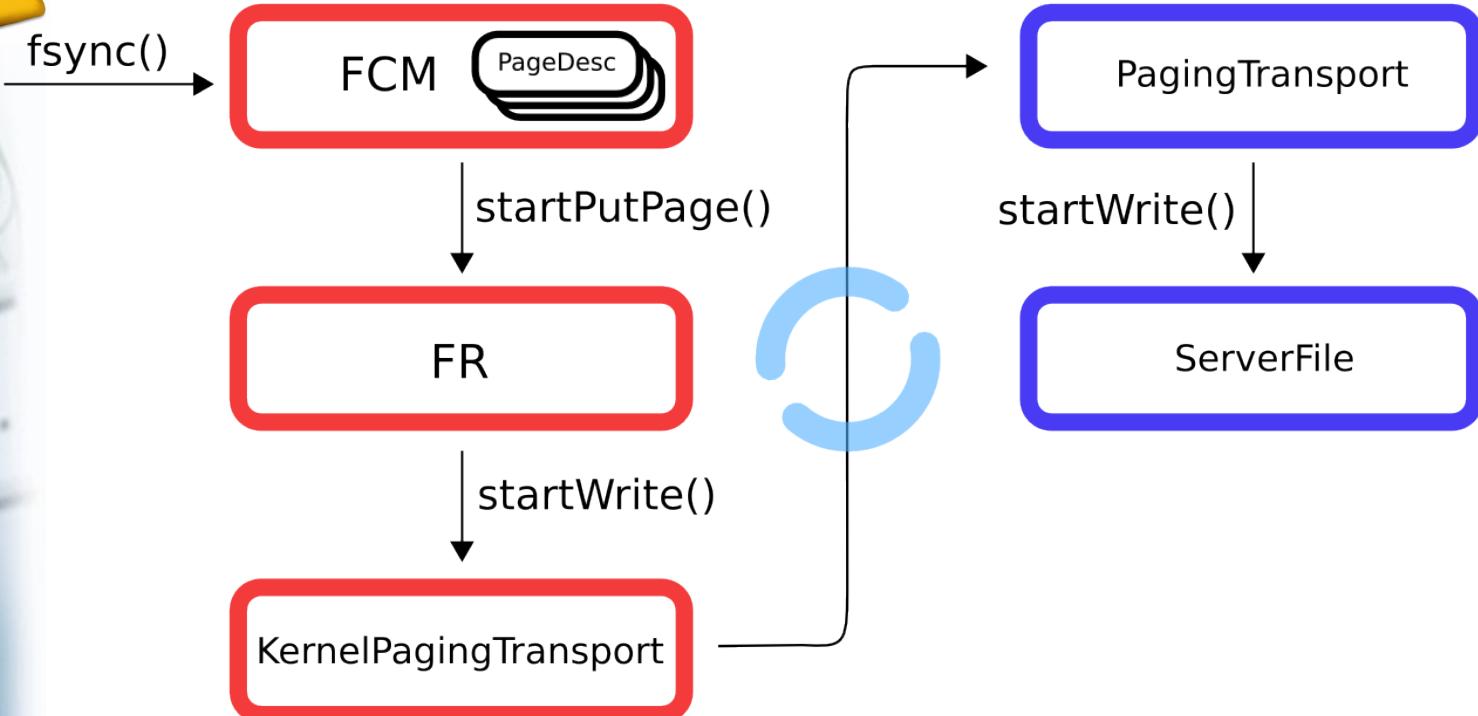
# Current fsync Path



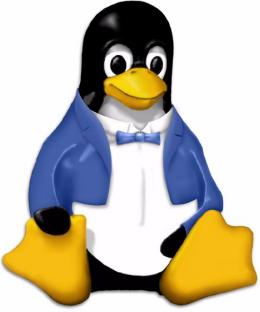
...over a shared ring-buffer, to the `PagingTransport` object



# Current fsync Path



ServerFile object handles write operation



# I/O Request Structure

- Basic unit of I/O
- Shared with userspace
- fileToken member
  - reference to ServerFile
  - opaque within kernel

PagingRequest

type

fileToken

fileOffset

addr

size



# K42 I/O Subsystem

- In summary:
  - FCM (File Cache Manager)
    - Manages mapped regions
  - FR (File Representative)
    - Holds references to userspace file objects
  - ServerFile
    - Userspace file object
  - I/O performed by sending requests to userspace
    - Requests are “pushed” to the device

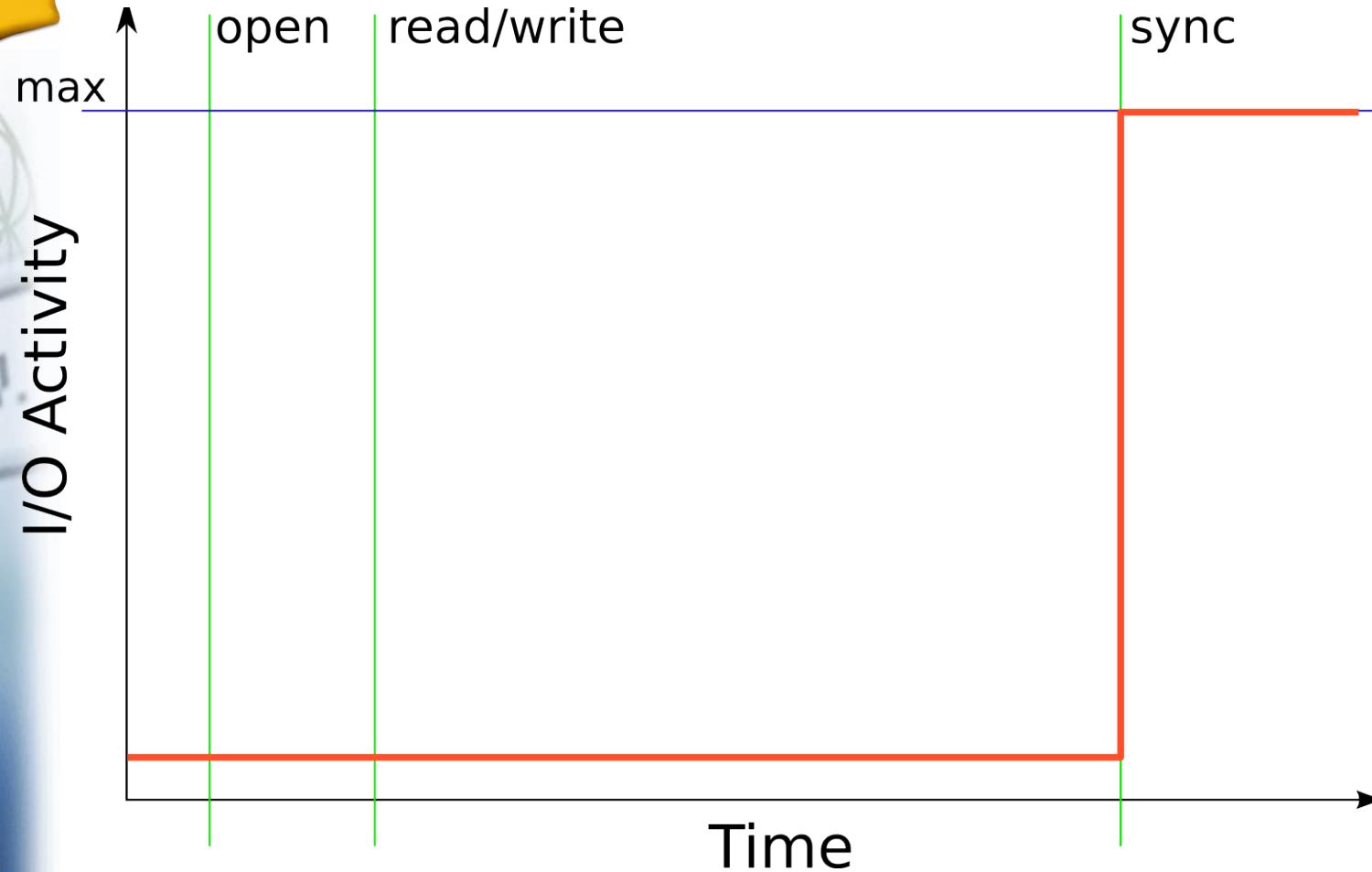


# Device-driven I/O

- Goal: Maximise usage of device bandwidth
  - Perform implicit operations when device is idle
- New programming model for I/O operations
  - FCMs no longer asynchronously generate implicit requests
  - Instead, FCMs are queried for requests
    - “pull model”

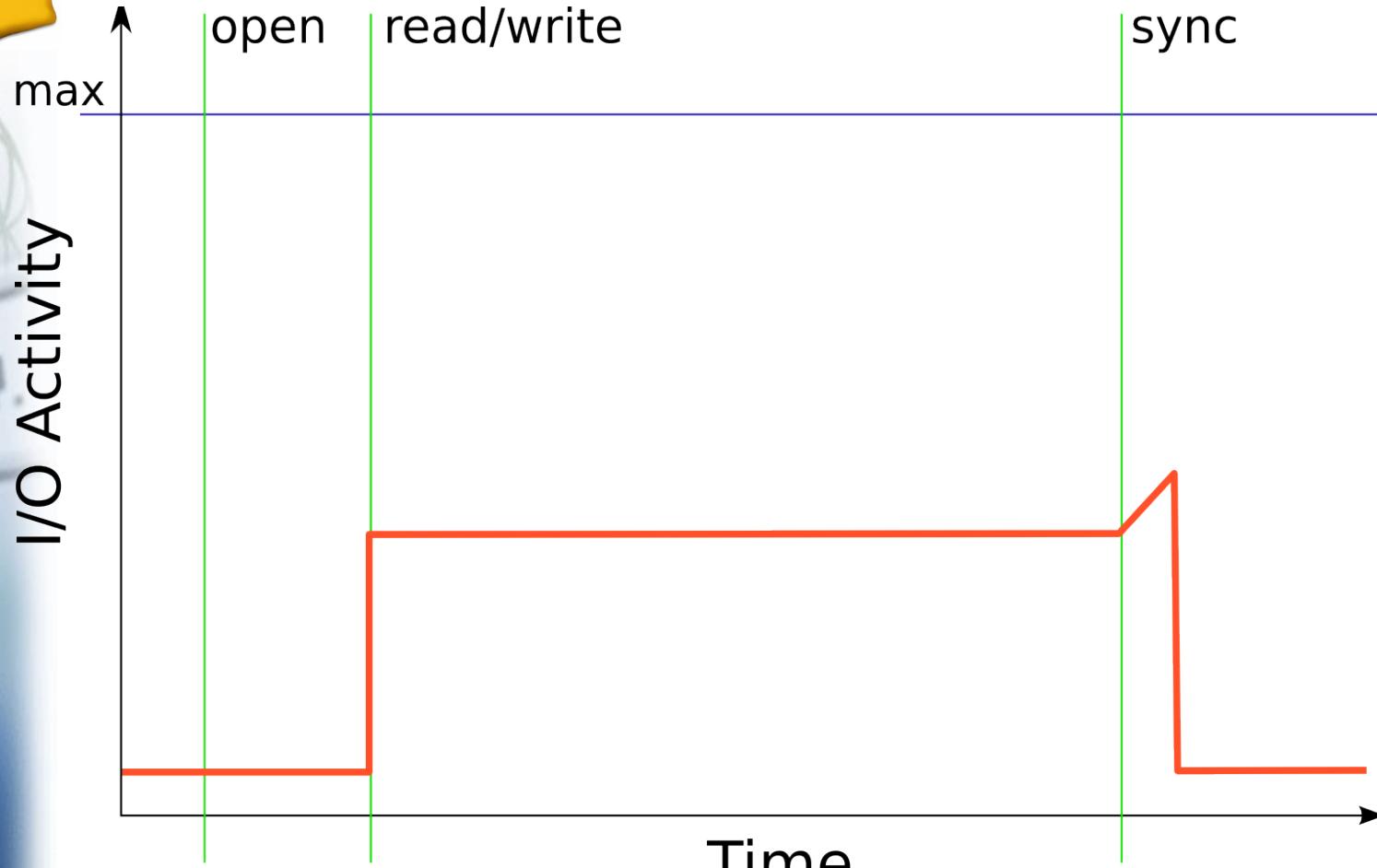


# Current I/O Activity





# Desired I/O Activity

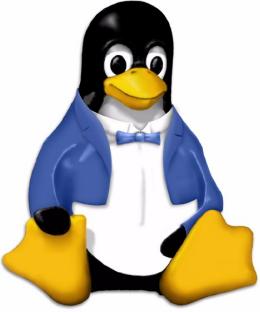


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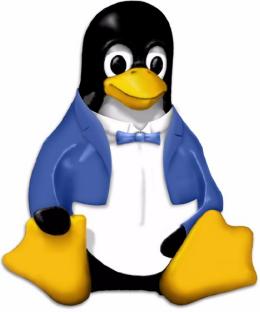
# Rationale

- Allows aggressive prefetch & page cleaning
  - Should stabilise bursty I/O
  - No need to be conservative
- FCMs “know best” how to optimise implicit ops
- Low request latency
  - Requests are only generated when the device is ready
  - No need for a cancellation mechanism



# Design

- Keep track of pageable FCMs
  - New class: PagingService
- Generate requests
  - New method: FCM::getRequests()
- Notify the transport that bandwidth is available
  - New PPC: KernelPagingTransport::notify()



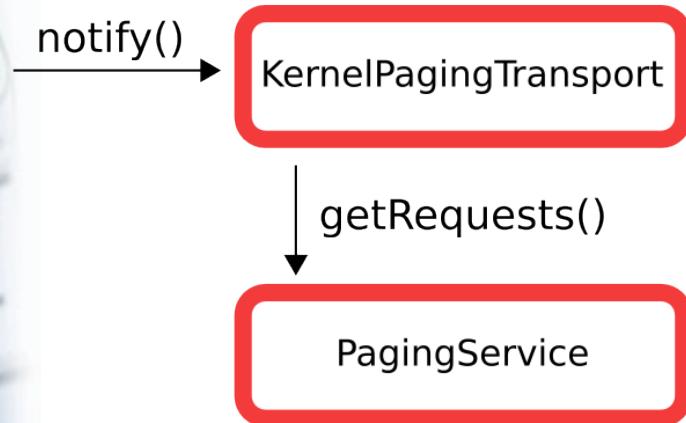
# Device-driven Design



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# Device-driven Design





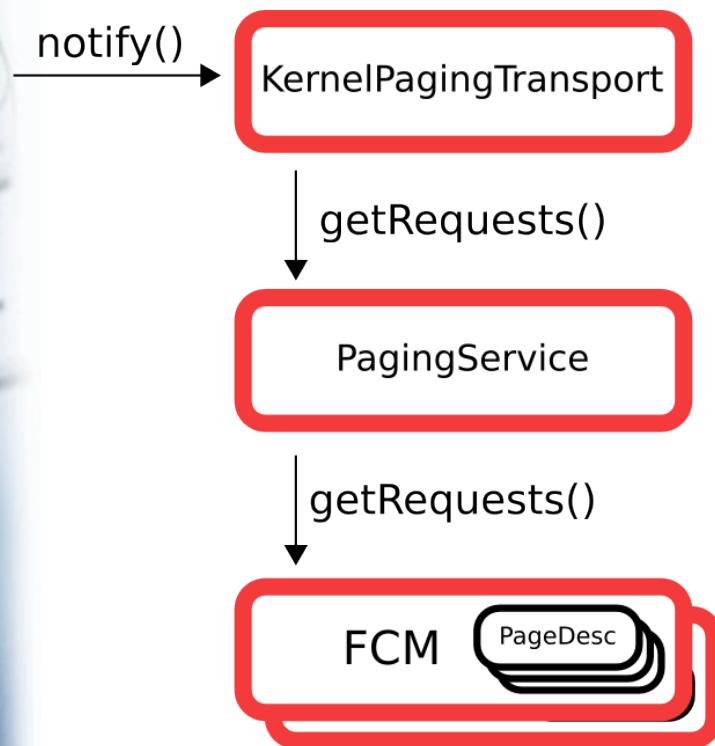
# The PagingService

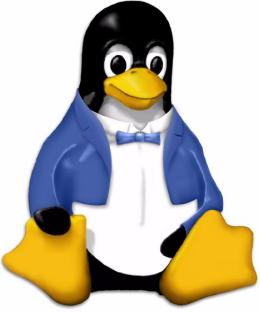
- Clustered object
  - One root object created at boot time
  - One “representative” created per processor
- FCMs register with the PagingService on creation
  - Only pageable FCMs
  - De-register on destruction
- Referenced by the KernelPagingTransport





# Device-driven Design



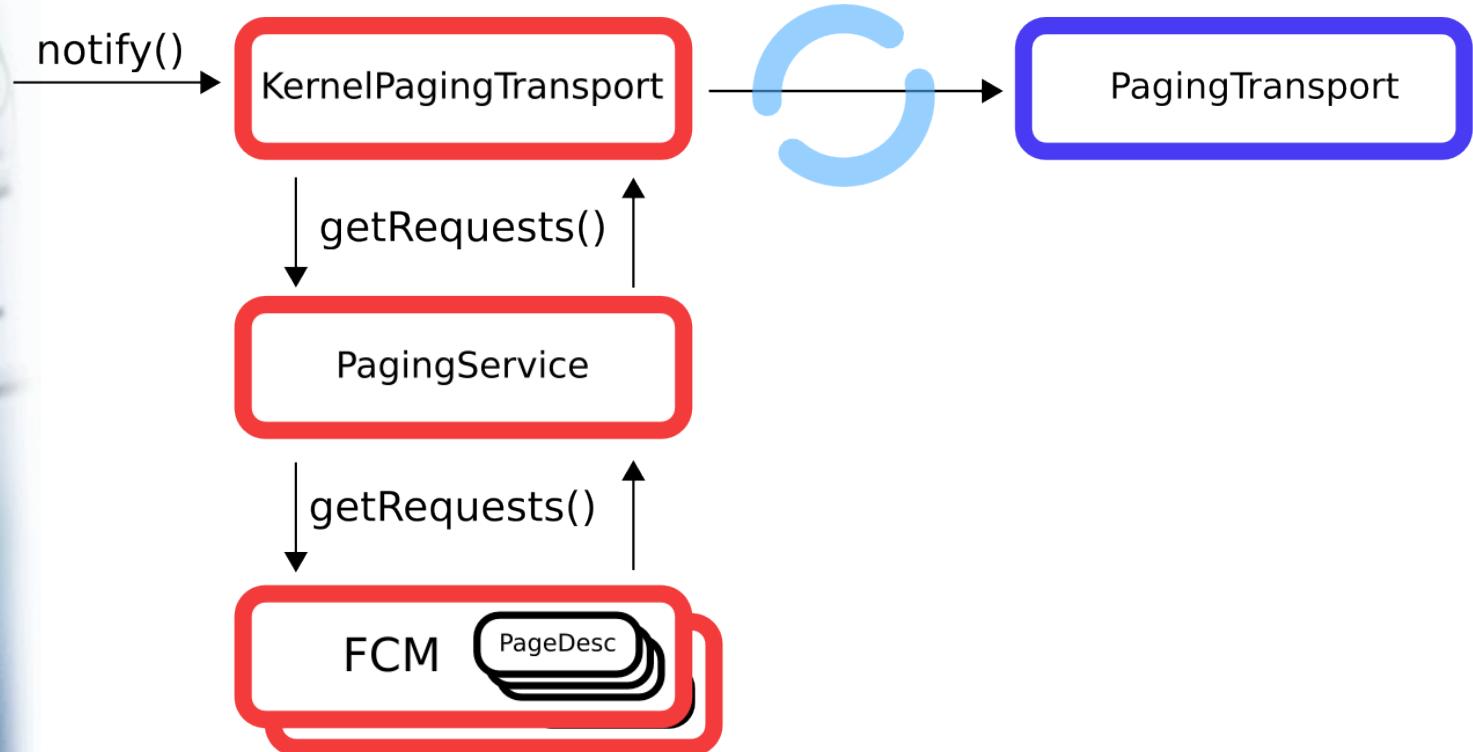


# FCM::getRequests()

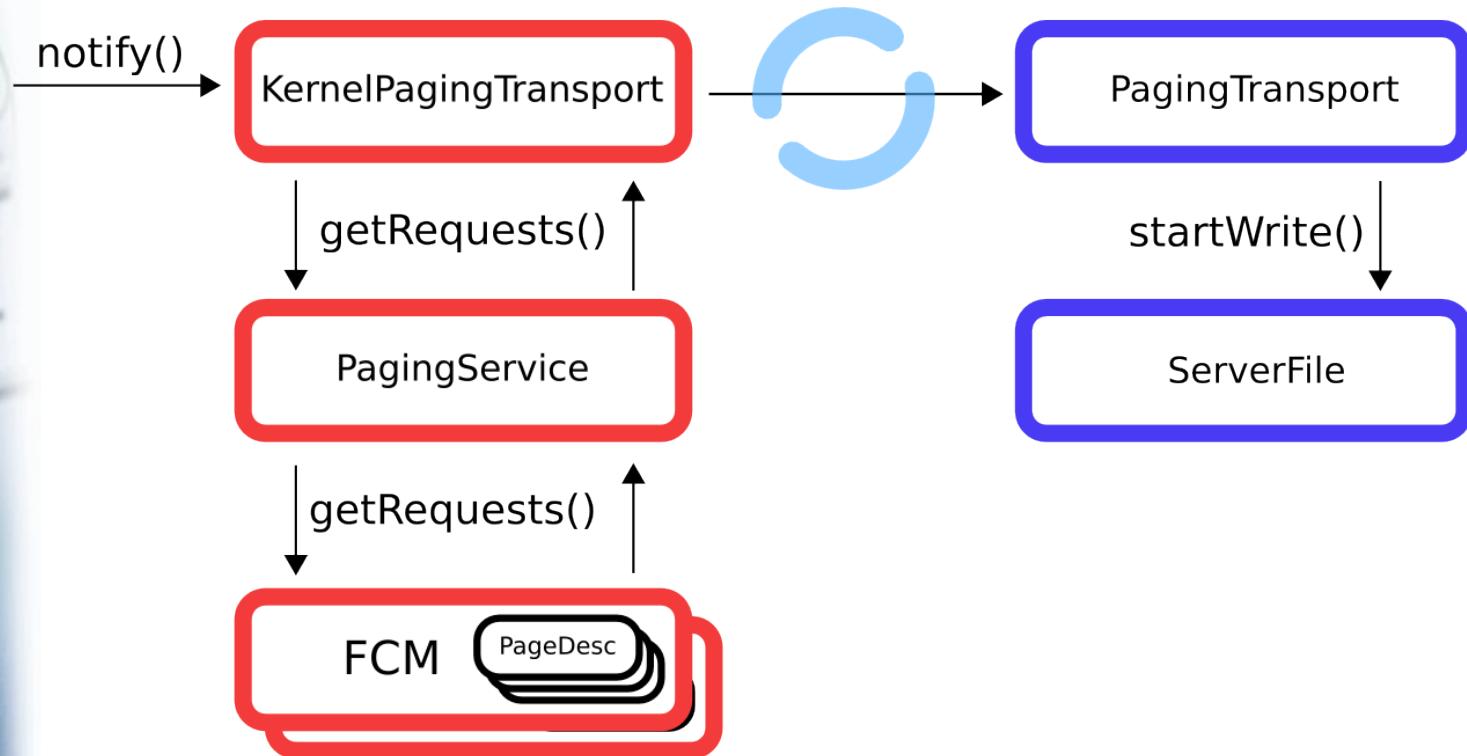
- Method to retrieve a set of I/O requests
  - Up to a specified maximum
- Implicit request pattern controlled by FCM
  - Can be based on history, expected patterns
    - Prefetch optimisation
    - Resident-set research at UofT
  - Easy to explore new ideas
    - FCM inheritance
  - Dynamic upgrade



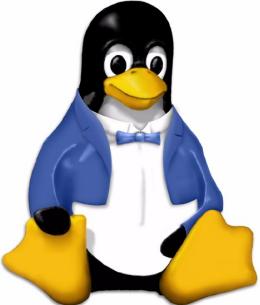
# Device-driven Design



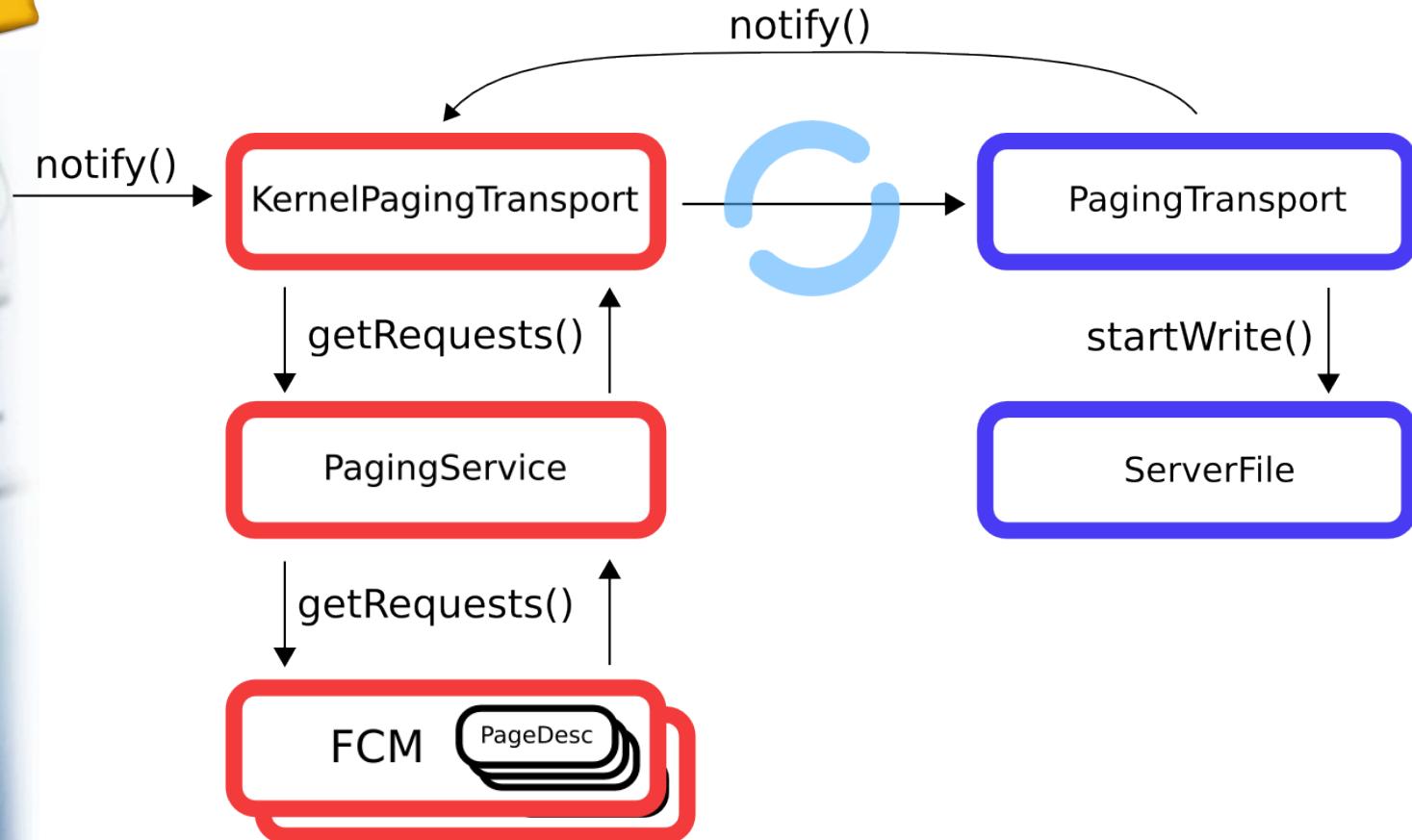
# Device-driven Design



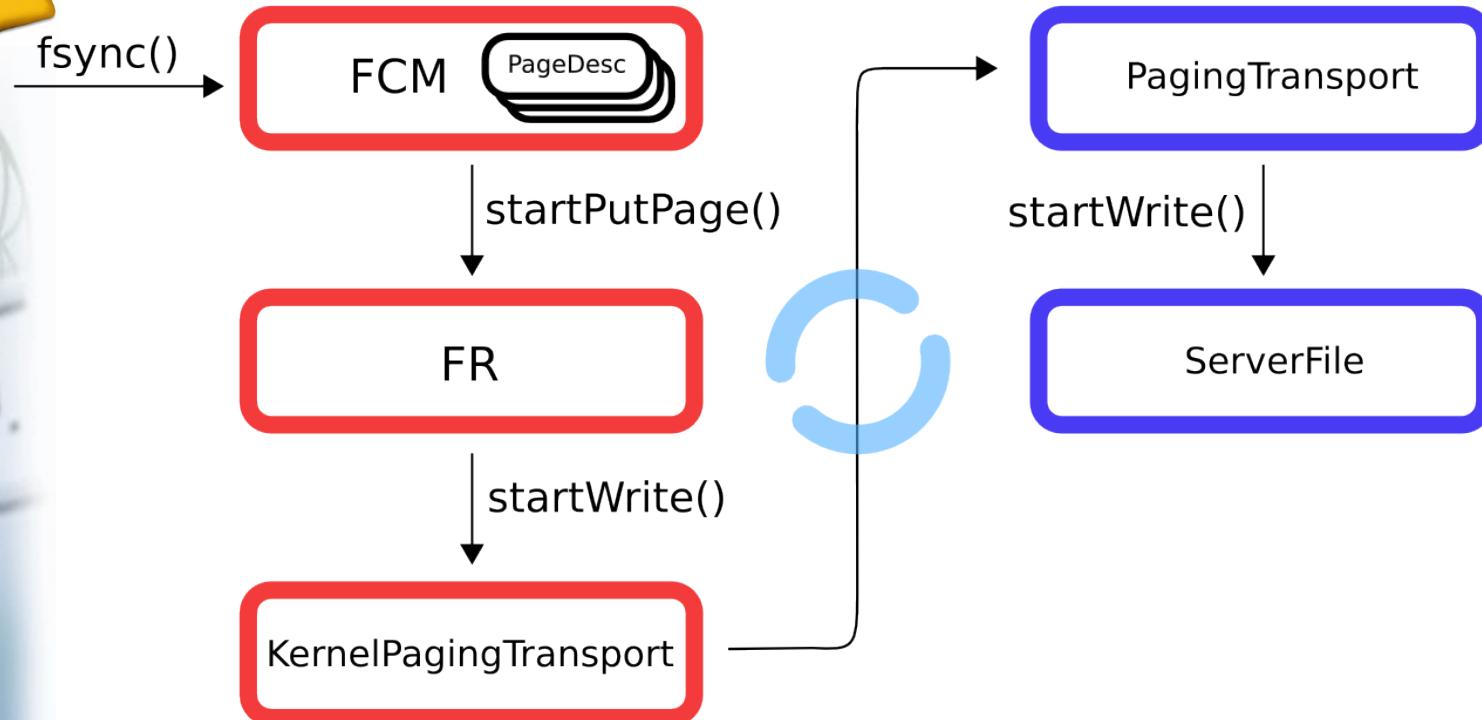
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# Device-driven Design

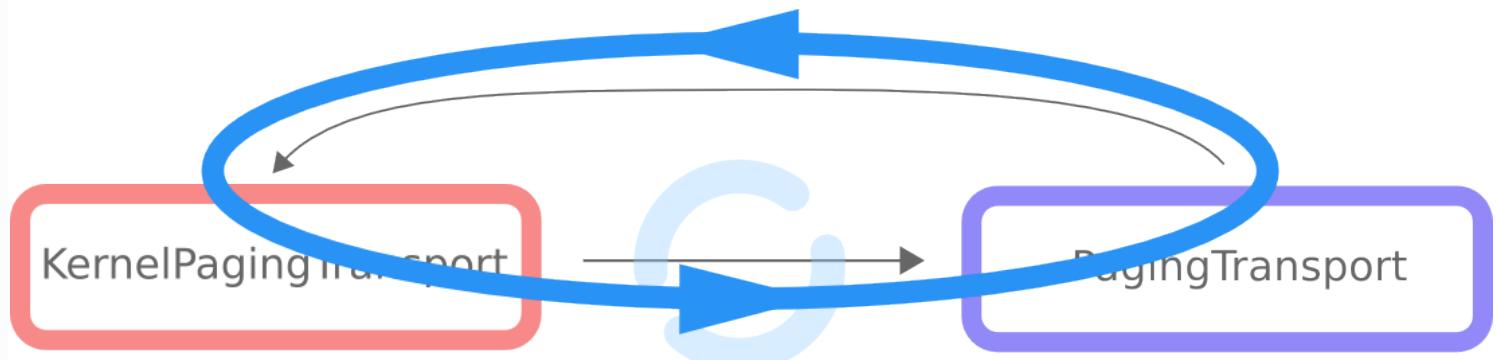


# Design: Current Model

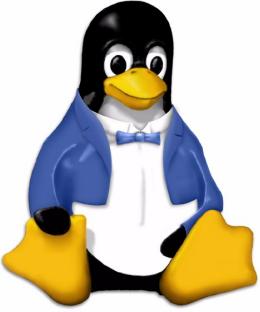




# Flow Control



- Feedback loop provides flow control
  - Restricted by ring-buffer availability
  - Representative of device usage



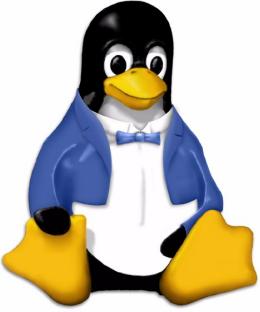
# Initial Concerns

- Global optimisation of implicit operations
  - No longer possible?
- Resource usage by PagingService
  - Don't want to tie up other resources
- Latency introduced for explicit operations



# Tunables

- Implementation of FCM::getRequests()
  - Prefetch optimisation, page cleaning patterns
  - A whole area of research in itself...
- Allocation of bandwidth to FCMs
  - Which to query for requests?
  - Priority between FCMs



# The Implementation

lib/libc/fslib/PagingTransport.C		13	+-
lib/libc/fslib/PagingTransport.H		11	+
lib/libc/io/FileLinux.C		8	-
os/kernel/ObjectRefsKern.H		4	
os/kernel/bilge/CObjGlobalsKern.H		8	-
os/kernel/init/KernelInit.C		4	
os/kernel/mem/FCM.C		6	+
os/kernel/mem/FCM.H		2	
os/kernel/mem/FCMDefaultMultiRep.C		4	
os/kernel/mem/FCMFile.C		55	++++++-
os/kernel/mem/FCMFile.H		6	+
os/kernel/mem/FR.H		11	+
os/kernel/mem/FRCommon.H		6	+
os/kernel/mem/FRPA.C		30	++++
os/kernel/mem/FRPA.H		3	
os/kernel/mem/KernelPagingTransport.C		60	++++++-
os/kernel/mem/KernelPagingTransport.H		16	+-
os/kernel/mem/Makefile		7	-
os/kernel/mem/PagingService.C		204	+++++++++++++++++++++
os/kernel/mem/PagingService.H		166	+++++++++++++++++++++
20 files changed, 585 insertions(+), 39 deletions(-)			



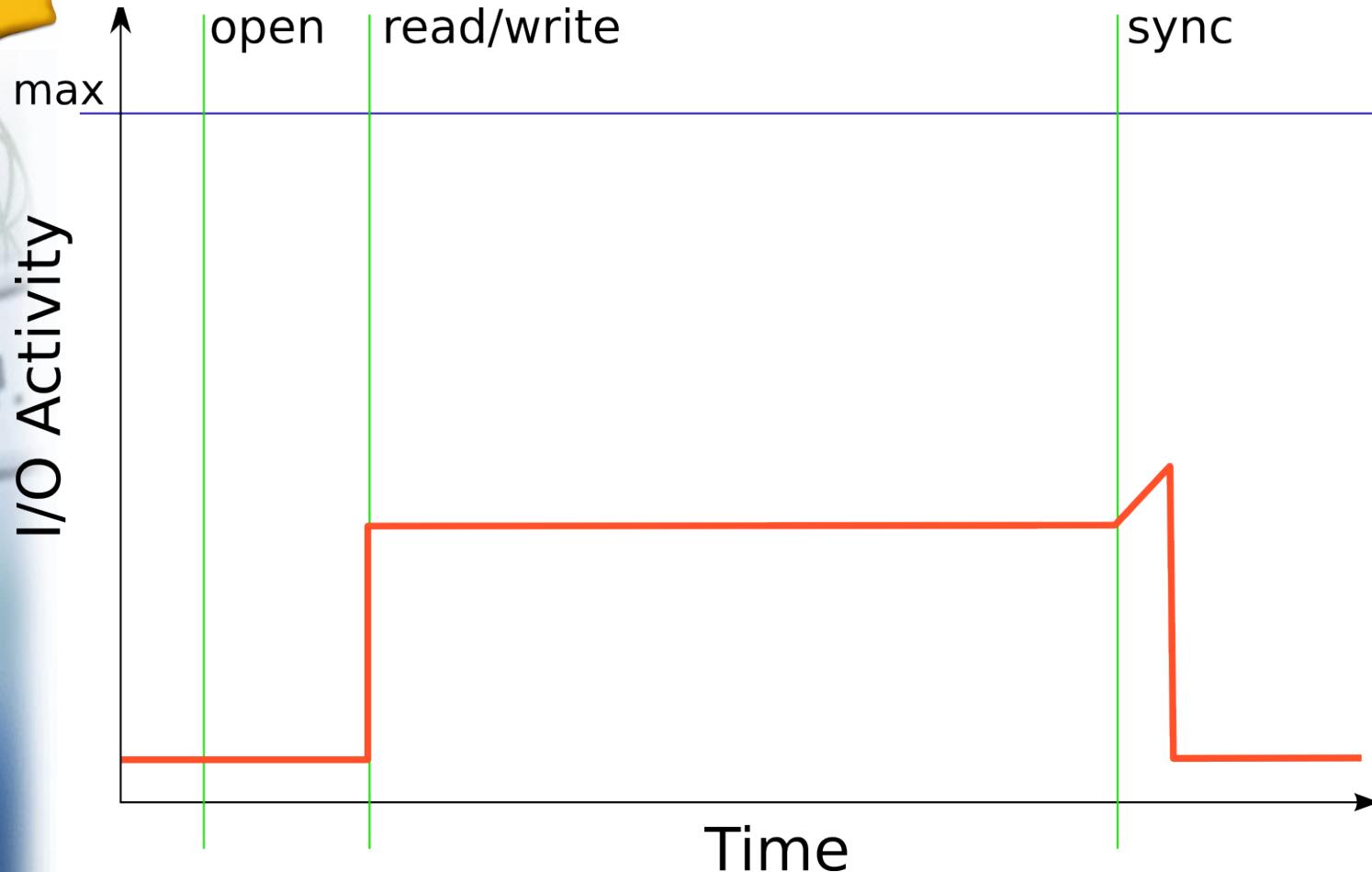


# Experimentation

- Initial testing
  - Copy between filesystems
  - Time taken to sync changes to disk
- Implementing a prefetch oracle
  - Advance knowledge of pattern
- Alter prefetch oracle
  - Use history to determine pattern



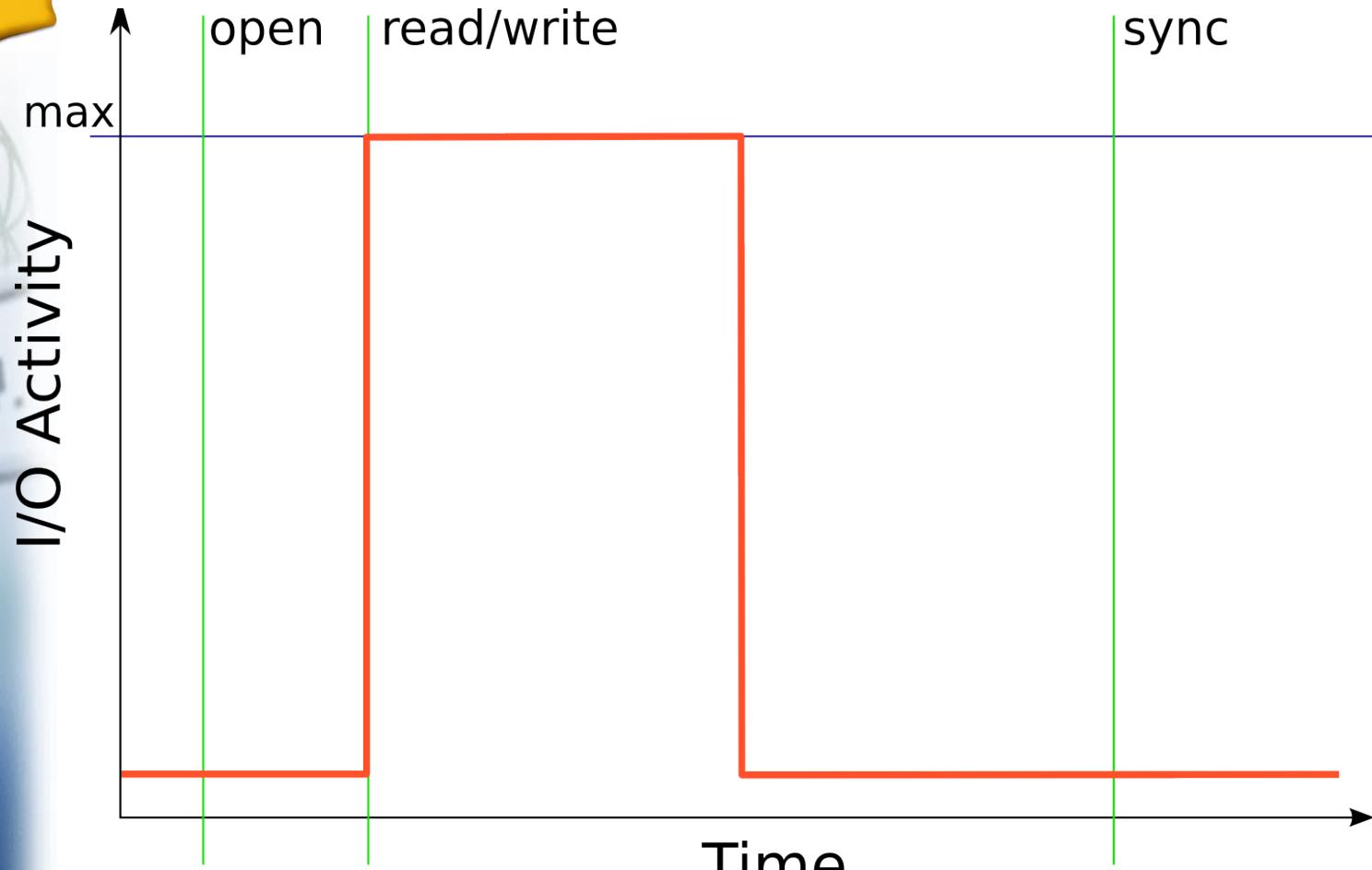
# Desired I/O Activity



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# Actual I/O Activity

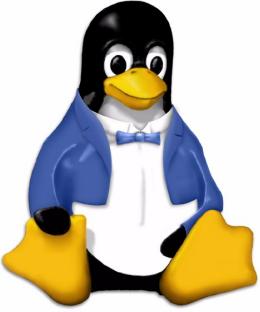


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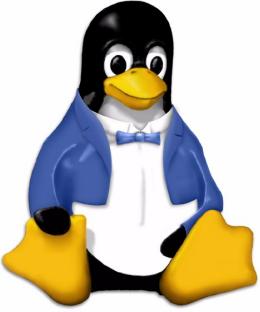
# Current Issues

- Current implementation is very aggressive
  - PagingService blocks other operations while paging
    - sync has no pages to flush
- High use of the PPC mechanism
  - Potential bottleneck



# Further Work

- Tune scheduler to allow true background paging
- Improve FCM-request allocation
- Implementation of prefetch/cleaning algorithm
  - FCM::getRequests()



# Resources

- K42 website at IBM Research
  - <http://www.research.ibm.com/K42/>
- Patches
  - <http://ozlabs.org/~jk/projects/k42/>



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